

# Introduction to Z3

h4tum Presentation

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- ▶ What is a **SMT** solver (and therefore **Z3**)?
- ▶ How to use **Z3**
- ▶ Hands-On Examples from real CTF challenges

# What is a SMT Solver?

- ▶ Satisfiability Modulo Theories is like SAT, but with non-boolean values
- ▶ Encompasses most mathematical formulas we can think of.
- ▶ SMT solvers try to solve the following problems:
  - ▶ Find a solution to a constraint system
  - ▶ Find the maximum/minimum of a formula in a constraint system
  - ▶ Simplify formulas
- ▶ Keep in mind: These problems are NP-hard in general.
- ▶ These solvers can solve many problems, but often not optimal (in terms of performance)

# What is Z3

- ▶ Z3 is a “theorem prover” by Microsoft Research
- ▶ Bundles several solvers including SAT and SMT
- ▶ State-of-the-art tool used by security researchers and CTF-players
- ▶ Can be relevant for any CTF category
- ▶ Includes bindings for 10 programming languages.
- ▶ Python API is easy-to-use. `import z3`

# How to use Z3? - The Solver object

- ▶ The base object for solving is the Solver object.
- ▶ Create solver: `s = z3.Solver()`
- ▶ Add constraints: `s.add(constraint)`
- ▶ Check for satisfiability: `s.check()`
- ▶ Get model (after check): `m = s.model()`
- ▶ Query model: `m[var].as_long()`

- ▶ Creating variables:
  - ▶ Arbitrary-precision integer: `z3.Int("v")`
  - ▶ Bitvector: `z3.BitVec("v", bitsize)`
- ▶ Creating formulas:
  - ▶ With normal python operators: `v * 0x42 == 1`  
**Be aware:** Normal operators are signed versions
  - ▶ For unsigned: `z3.{ULT, ULE, UGT, UGE, UDiv, URem, LShR}`
  - ▶ Further logical operators:
    - ▶ If-Then-Else: `z3.If(cond, then, else)`
    - ▶ Variables must be different: `z3.Distinct(list_of_vars)`
    - ▶  $A \implies B$ : `z3.Implies(A, B)`

- ▶ **Bitvectors** are big-endian
- ▶ Mix up signed and unsigned operations (especially bitshifts)
- ▶ Query **model** before **sat** check

- ▶ **crypto** Challenge from HTB Cyber Apocalypse CTF 2024
- ▶ Official difficulty: **insane**
- ▶ We are asked to solve the following problem:

Given  $h1$  and  $h2$ , find  $s1, s2, r1, r2, r3, r4$  s.t.:

$$\text{ROL}(s1, r1) \oplus \text{ROL}(s2, r2) == h1$$

$$\text{ROL}(s1, r3) \oplus \text{ROL}(s2, r4) == h2$$

- ▶ Using **some** mathematical reasoning, we can solve this if we can solve:

Given  $x$ , find  $y, c$  s.t.:

$$y \oplus \text{ROL}(y, c) == x$$

- ▶ **Task**: Find a solution with **Z3** - template code given

- ▶ `pwn/misc` Challenge from GPN CTF 2024
- ▶ Points: 313
- ▶ We can execute 4 bytes of arbitrary shellcode followed by 100 bytes of shellcode generated via a bad **PRNG**
- ▶ We created a payload with 9 bytes of shellcode  
→ we need to make the **PRNG** generate 5 specific bytes
- ▶ **Task**: Find a solution with **Z3** - template code given
  - ▶ `test.py` for testing the implementation
  - ▶ `exploit.py` to get the flag
  - ▶ Implement the function **`get_seed_for_sequence`**

- ▶ **rev** Challenge from GPN CTF 2024
- ▶ Points: 135
- ▶ License key checker with basically 5 binaries of different architecture (main + 4 sub-binaries)
- ▶ Each binary checks some constraints on the license key
- ▶ Reversing is already performed and text-descriptions given
- ▶ License is a string of form **XXXXX-XXXXX-XXXXX-XXXXX** where **Xs** are variables
- ▶ **Task**: Model all constraints with **Z3** and retrieve the license key
- ▶ Note: template code given again